Homotopy height, grid-major height and graph-drawing height

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#### Problem statement

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# Applications

Drawing planar graphs on narrow strips of paper



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Measuring similarity between curves on surfaces



All our graphs are planar



All our graphs are planar All faces (including the outer face) are triangular



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Models a triangulated sphere























How short of a curve can sweep a topological sphere? Variant in this talk: curve fixed to arbitrary basepoint Homotopy height =  $\inf_{basepoint} \inf_{sweep} \sup_t \|sweep(t)\|$ 



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Triangulate surface to approximate metric



Triangulate surface to approximate metric

Basepoint = face of triangulation



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Basepoint = face of triangulation = outer face



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Basepoint = face of triangulation = outer face All curves  $\gamma_t$  of sweep start and end on outer face



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Basepoint = face of triangulation = outer face All curves  $\gamma_t$  of sweep start and end on outer face First and last curves of sweep consist of single (distinct) vertex



Triangulate surface to approximate metric

Basepoint = face of triangulation = outer face All curves  $\gamma_t$  of sweep start and end on outer face First and last curves of sweep consist of single (distinct) vertex Consecutive curves differ by a (simple) homotopy move



# Simple homotopy moves

Any curve in simple sweep uses any vertex  $\leq$  once

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Homotopy moves (nonsimple)

Vertices can be reused

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Simple homotopy moves + edge spikes:



Homotopy moves (nonsimple)

Vertices can be reused

Simple homotopy moves + edge spikes:



Sweep must flip (or slide) across each face 'from-left-to-right' once more than 'from-right-to-left'



 $\begin{array}{l} W \!\!\times \!\! H \text{ gridpoints} \\ \{1, \dots, W\} \times \{1, \dots, H\} \end{array}$ 







 $W \!\!\times\! H$  grid

graph on gridpoints, edges between points at distance 1





#### $W \!\!\times\! H$ grid

graph on gridpoints, edges between points at distance 1

Grid-major height (of a planar graph G) minimum h s.t. G is a minor of  $W \times h$  grid




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Grid-major height (of a planar graph G) minimum h s.t. G is a minor of  $W \times h$  grid

 $\begin{array}{l} \mbox{Minor (of graph $H$)} \\ \mbox{graph obtained from $H$ by} \\ \mbox{contracting edges} \\ \mbox{removing edges/vertices} \end{array}$ 





#### $W \times H$ grid

graph on gridpoints, edges between points at distance  $1 \$ 

Grid-major height (of a planar graph G) minimum h s.t. G is a minor of  $W \times h$  grid

Minor (of graph H) graph obtained from H by contracting edges removing edges/vertices Simple grid-major height each label in a column appears consecutively Some graph parameters...

(Simple) homotopy height

(Simple) grid-major height

(Simple) contact representation height

Visibility representation height

Straight-line drawing height

Pathwidth

Outerplanarity

Contact representation each gridpoint labeled by a vertex of G





each gridpoint labeled by a vertex of Geach label forms connected subgraph two labels adjacent if and only if edge in G





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Simple contact representation

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### Simple contact representation

each label appears consecutively in each column

#### (Simple) contact representation height

min h s.t.  $W \times h$  grid has (simple) contact representation





each vertex corresponds to a horizontal bar





each vertex corresponds to a horizontal bar for each edge there is a line of visibility (horizontal or vertical)







each vertex corresponds to a horizontal bar for each edge there is a line of visibility (horizontal or vertical) bars and lines of visibility do not cross



Visibility representation height min h s.t.  $W \times h$  grid has flat visibility representation











Straight-line height

min h with planar straight line drawing that has all vertices on  $W \times h$  gridpoints



Straight-line height

min h with planar straight line drawing that has all vertices on  $W \times h$  gridpoints



# Outerplanarity

#### Outerplanarity (of a planar embedding) number of steps needed to remove all vertices each step: remove vertices of outer face



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#### Outerplanarity (of a planar embedding) number of steps needed to remove all vertices each step: remove vertices of outer face

Outerplanarity (of a planar graph) minimum outerplanarity over all embeddings

# Pathwidth

#### Path decomposition

Form groups of vertices and put groups on a path Each vertex belongs to a subpath of groups For any edge, endpoints lie in a common group



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#### Path decomposition

Form groups of vertices and put groups on a path Each vertex belongs to a subpath of groups For any edge, endpoints lie in a common group

#### Pathwidth

Minimum largest group size -1 over all decompositions



Relations between graph parameters...

(Simple) homotopy height

(Simple) grid-major height

(Simple) contact representation height

Visibility representation height

Straight-line drawing height

Pathwidth

Outerplanarity

Every contact representation is a grid-major representation



Every contact representation is a grid-major representation

Reverse is not necessarily true:

Grid-major repr. can have unwanted contacts and empty spots



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Our assumptions on the graph

 $\Rightarrow$  empty space can be filled without unwanted contacts



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Our assumptions on the graph  $\Rightarrow$  empty space can be filled without unwanted contacts

contact representation height = grid-major height

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Our assumptions on the graph  $\Rightarrow$  empty space can be filled without unwanted contacts

contact representation height = grid-major height

simple contact representation height = simple grid-major height

Every contact representation is a grid-major representation

Reverse is not necessarily true: Grid-major repr. can have unwanted contacts and empty spots

Our assumptions on the graph  $\Rightarrow$  empty space can be filled without unwanted contacts

contact representation height = grid-major height

simple contact representation height = simple grid-major height

Requiring that regions are x-monotone can only increase height

grid-major height  $\leq$  simple grid-major height

Every flat visibility representation can be turned into a simple grid-major representation



Every flat visibility representation can be turned into a simple grid-major representation

simple grid-major height  $\leq$  visibility representation height



Every flat visibility representation can be turned into a simple grid-major representation

simple grid-major height  $\leq$  visibility representation height Previously shown [Biedl14]:

visibility representation height = straight-line drawing height



Pathwidth of  $W\!\mathbf{x}h$  grid minor  $\leq$  pathwidth of  $W\!\mathbf{x}h$  grid  $\leq h$ 

Pathwidth of Wxh grid minor  $\leq$  pathwidth of Wxh grid  $\leq h$ pathwidth  $\leq$  grid-major height

Pathwidth of W x h grid minor  $\leq$  pathwidth of W x h grid  $\leq h$ pathwidth  $\leq$  grid-major height

Outerplanarity of  $W \times h$  grid minor  $\leq$  that of  $W \times h$  grid  $\leq \lceil h/2 \rceil$ 

 $2 \text{ outerplanarity } -1 \leq \operatorname{grid-major}$  height

# Overview of bounds

```
2 \text{ outerplanarity } -1 \text{ and } \text{ pathwidth}
          grid-major height
    contact representation height
      simple grid-major height
simple contact representation height
   visibility representation height
     straight-line drawing height
```



Simple grid-major height = simple homotopy height

Sweep can be assumed monotone based on [CMO et al. 17] curve does not sweep backwards

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Simple homotopy height  $\geq$  simple grid-major height:


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Simple homotopy height  $\leq$  simple grid-major height:

Take contact representation

wlog 3 colors on boundary



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Simple homotopy height  $\leq$  simple grid-major height:

Take contact representation wlog 3 colors on boundary No four polygons meet at a point





Sweep can be assumed monotone based on [CMO et al. 17] curve does not sweep backwards

Simple homotopy height  $\leq$  simple grid-major height:

or

Take contact representation wlog 3 colors on boundary No four polygons meet at a point

Remove interior vertical junctions



Sweep can be assumed monotone based on [CMO et al. 17] curve does not sweep backwards

Simple homotopy height  $\leq$  simple grid-major height:

Take contact representation wlog 3 colors on boundary No four polygons meet at a point

Remove interior vertical junctions

$$\rightarrow$$
 or



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Simple homotopy height  $\leq$  simple grid-major height:

Take contact representation wlog 3 colors on boundary No four polygons meet at a point

Remove interior vertical junctions

$$\rightarrow$$

Make *x*-coordinates distinct



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Simple homotopy height  $\leq$  simple grid-major height:

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 or

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Sweep can be assumed monotone based on [CMO et al. 17] curve does not sweep backwards

Simple homotopy height  $\leq$  simple grid-major height:

Take contact representation wlog 3 colors on boundary No four polygons meet at a point

Remove interior vertical junctions





Make x-coordinates distinct Make left and right boundary single (but distinct) color

Sweep can be assumed monotone based on [CMO et al. 17] curve does not sweep backwards





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Simple homotopy height  $\leq$  simple grid-major height: Extract sweep





Similarly, grid-major height = homotopy height





# Overview of bounds 2 outerplanarity -1 and pathwidth grid-major height contact representation height = homotopy height simple grid-major height simple contact representation height = simple homotopy height visibility representation height inequalities are strict straight-line drawing height gaps are nonconstant













#### Pathwidth $\leq$ grid-major height

Pathwidth = 3 Grid-major height  $\geq 2$  outerplanarity  $-1 \geq n/3 - 1$ 


# Pathwidth $\leq$ grid-major height

# $\begin{array}{l} \mbox{Pathwidth}=3\\ \mbox{Grid-major height}\geq 2 \mbox{ outerplanarity } -1\geq n/3-1\\ n/6 \mbox{ triangles will be nested, no matter the outer face} \end{array}$



#### Outerplanarity $\leq$ grid-major height



Outerplanarity  $\leq$  grid-major height Grid-major height  $\geq$  pathwidth  $= \Omega(\log n)$ 



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Nonsimple  $\leq$  simple grid-major height









Minor of and hence of *W*x4 grid



Minor of  $\chi$  and hence of Wx4 grid  $\Rightarrow$  grid-major height  $\leq 4$ 





Grid-major height  $\leq 4$ Simple grid-major height  $= \Omega(n)$ :



Grid-major height  $\leq 4$ Simple grid-major height  $= \Omega(n)$ : Diameter of subgraph is  $\Omega(n)$ 



 $\begin{array}{l} \mbox{Grid-major height} \leq 4 \\ \mbox{Simple grid-major height} = \Omega(n): \\ \mbox{Diameter of subgraph is } \Omega(n) \\ \mbox{Some vertex in subgraph is far from 'outer face'} \end{array}$ 



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Grid-major height  $\leq 4$ 

Simple grid-major height =  $\Omega(n)$ :

Diameter of subgraph is  $\Omega(n)$ 

Some vertex in subgraph is far from 'outer face'

That vertex splits some path in sweep in two pieces



Grid-major height  $\leq 4$ 

Simple grid-major height =  $\Omega(n)$ :

Diameter of subgraph is  $\Omega(n)$ 

Some vertex in subgraph is far from 'outer face'

That vertex splits some path in sweep in two pieces

At least one piece lies in subgraph, and is therefore long

For series-parallel graphs, simple grid-major height is  $O(\log n)$ 

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For series-parallel graphs, simple grid-major height is  $O(\log n)$ Contact-representation with source/target in top/bottom-right



edge

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For series-parallel graphs, simple grid-major height is  $O(\log n)$ Contact-representation with source/target in top/bottom-right



series

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series

#### parallel

Height increases (by 2) only if combined grids are similar height  $\Rightarrow$  grid-major height =  $O(\log n)$ 

For series-parallel graphs, simple grid-major height is  $O(\log n)$ 

There exist series-parallel graphs with graph-drawing height =  $\Omega(2^{\sqrt{\log n}})$  [Frati10]

For series-parallel graphs, simple grid-major height is  $O(\log n)$ 

There exist series-parallel graphs with graph-drawing height =  $\Omega(2\sqrt{\log n})$  [Frati10]

Triangulating them cannot decrease height



