Homotopy height, grid-major height and graph-drawing height

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## Problem statement

Given a planar graph and a height $h$, is there a planar straight line drawing of height $h$ ?

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Applications
Drawing planar graphs on narrow strips of paper


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Measuring similarity between curves on surfaces


Assumptions on our graphs
All our graphs are planar


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All faces (including the outer face) are triangular


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$\Rightarrow$ choice of outer face fully determines rotation system


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All faces (including the outer face) are triangular
$\Rightarrow$ choice of outer face fully determines rotation system
Models a triangulated sphere


## Homotopy height

How short of a curve can sweep a topological sphere?

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## Discretizing Homotopy height

Triangulate surface to approximate metric


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## Discretizing Homotopy height

Triangulate surface to approximate metric
Basepoint $=$ face of triangulation $=$ outer face All curves $\gamma_{t}$ of sweep start and end on outer face
First and last curves of sweep consist of single (distinct) vertex Consecutive curves differ by a (simple) homotopy move


## Simple homotopy moves

Any curve in simple sweep uses any vertex $\leq$ once

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Boundary-move



Boundary-edge-slide


## Simple homotopy moves

Any curve in simple sweep uses any vertex $\leq$ once

Face-flip
(not outer face)

Boundary-move



Boundary-edge-slide


## Homotopy moves (nonsimple)

## Vertices can be reused

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Simple homotopy moves + edge spikes:

## Homotopy moves (nonsimple)

Vertices can be reused
Simple homotopy moves + edge spikes:

Sweep must flip (or slide) across each face 'from-left-to-right' once more than 'from-right-to-left'

Grid-major heightw!
$W \times H$ gridpoints $\{1, \ldots, W\} \times\{1, \ldots, H\}$

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W×H grid

graph on gridpoints, edges between points at distance 1

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WxH grid
 graph on gridpoints, edges between points at distance 1

Grid-major height (of a planar graph $G$ ) minimum $h$ s.t. $G$ is a minor of $W \times h$ grid

Grid-major height N !
W×H gridpoints $\{1, \ldots, W\} \times\{1, \ldots, H\}$

W×H grid graph on gridpoints, edges between points at distance 1

Grid-major height (of a planar graph $G$ ) minimum $h$ s.t. $G$ is a minor of $W \times h$ grid

Minor (of graph H)
graph obtained from $H$ by
contracting edges removing edges/vertices

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Minor (of graph $H$ )
graph obtained from $H$ by contracting edges removing edges/vertices

Simple grid-major height each label in a column appears consecutively

## Some graph parameters...

(Simple) homotopy height
(Simple) grid-major height
(Simple) contact representation height
Visibility representation height
Straight-line drawing height
Pathwidth
Outerplanarity

Contact representation
each gridpoint labeled by a vertex of $G$


## Contact representation

each gridpoint labeled by a vertex of $G$
each label forms connected subgraph two labels adjacent if and only if edge in $G$


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each label appears consecutively in each column
(Simple) contact representation height $\min h$ s.t. $W \times h$ grid has
(simple) contact representation


Flat visibility representation
each vertex corresponds to a horizontal bar


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each vertex corresponds to a horizontal bar for each edge there is a line of visibility
(horizontal or vertical) bars and lines of visibility do not cross


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## Visibility representation height

 $\min h$ s.t. $W \times h$ grid has flat visibility representation

Flat visibility representation
each vertex corresponds to a horizontal bar for each edge there is a line of visibility
(horizontal or vertical)
bars and lines of visibility do not cross we allow additional visibilities (without edge in $G$ )
Visibility representation height
$\min h$ s.t. $W \times h$ grid has flat visibility representation


## Straight-line height

 $\min h$ with planar straight line drawing that has all vertices on $W \times h$ gridpoints

## Straight-line height

$\min h$ with planar straight line drawing that has all vertices on $W \times h$ gridpoints


## Outerplanarity

Outerplanarity (of a planar embedding) number of steps needed to remove all vertices each step: remove vertices of outer face

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Outerplanarity (of a planar embedding) number of steps needed to remove all vertices each step: remove vertices of outer face

Outerplanarity (of a planar graph) minimum outerplanarity over all embeddings

## Pathwidth

## Path decomposition

Form groups of vertices and put groups on a path
Each vertex belongs to a subpath of groups
For any edge, endpoints lie in a common group


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## Pathwidth

Minimum largest group size -1 over all decompositions


## Relations between graph parameters...

(Simple) homotopy height
(Simple) grid-major height
(Simple) contact representation height
Visibility representation height
Straight-line drawing height
Pathwidth
Outerplanarity

## Bounds

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Grid-major repr. can have unwanted contacts and empty spots


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Our assumptions on the graph
$\Rightarrow$ empty space can be filled without unwanted contacts
contact representation height $=$ grid-major height
simple contact representation height $=$ simple grid-major height
Requiring that regions are $x$-monotone can only increase height grid-major height $\leq$ simple grid-major height

## Bounds

Every flat visibility representation can be turned into a simple grid-major representation


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simple grid-major height $\leq$ visibility representation height


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Every flat visibility representation can be turned into a simple grid-major representation
simple grid-major height $\leq$ visibility representation height
Previously shown [Biedl14]:
visibility representation height $=$ straight-line drawing height


## Bounds

Pathwidth of $W \times h$ grid minor $\leq$ pathwidth of $W \times h$ grid $\leq h$

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Pathwidth of $W \times h$ grid minor $\leq$ pathwidth of $W \times h$ grid $\leq h$ pathwidth $\leq$ grid-major height
Outerplanarity of $W \times h$ grid minor $\leq$ that of $W \times h$ grid $\leq\lceil h / 2\rceil$
2 outerplanarity $-1 \leq$ grid-major height

## Overview of bounds

2 outerplanarity -1 and pathwidth

$$
\begin{gathered}
\quad \leq \\
\text { grid-major height } \\
=
\end{gathered}
$$

contact representation height

$$
\stackrel{\leq}{\leq}
$$

simple contact representation height

$$
\leq
$$

visibility representation height
straight-line drawing height

## Overview of bounds

2 outerplanarity -1 and pathwidth
$\qquad$
grid-major height
$=$
contact representation height $=$ homotopy height
$\leq \quad \leq$
simple contact representation height $=$ simple homotopy height
$\qquad$
visibility representation height
straight-line drawing height

## Simple grid-major height $=$ simple homotopy height

Sweep can be assumed monotone based on [CMO et al. 17] curve does not sweep backwards

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 Sweep can be assumed monotone based on [CMO et al. 17] curve does not sweep backwardsSimple homotopy height $\geq$ simple grid-major height:


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Simple homotopy height $\leq$ simple grid-major height:
Take contact representation wlog 3 colors on boundary


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Simple homotopy height $\leq$ simple grid-major height: Take contact representation wlog 3 colors on boundary No four polygons meet at a point


Simple grid-major height $=$ simple homotopy height
Sweep can be assumed monotone based on [CMO et al. 17] curve does not sweep backwards

Simple homotopy height $\leq$ simple grid-major height: Take contact representation wlog 3 colors on boundary No four polygons meet at a point x


Remove interior vertical junctions

$$
\square \rightarrow \square \text { or } \square
$$



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Make $x$-coordinates distinct

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Simple homotopy height $\leq$ simple grid-major height: Take contact representation wlog 3 colors on boundary No four polygons meet at a point


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\square \rightarrow \square
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Make $x$-coordinates distinct
Make left and right boundary single (but distinct) color

## Simple grid-major height $=$ simple homotopy height

Sweep can be assumed monotone based on [CMO et al. 17] curve does not sweep backwards
Simple homotopy height $\leq$ simple grid-major height: Extract sweep


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Simple homotopy height $\leq$ simple grid-major height: Extract sweep


Similarly, grid-major height $=$ homotopy height

## Overview of bounds

2 outerplanarity -1 and pathwidth
$\qquad$
grid-major height
$=$
contact representation height $=$ homotopy height
$\leq \quad \leq$
simple contact representation height $=$ simple homotopy height
$\qquad$
visibility representation height
straight-line drawing height

## Overview of bounds

2 outerplanarity -1 and pathwidth
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$$
=
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inequalities are strict
straight-line drawing height

## Overview of bounds

2 outerplanarity -1 and pathwidth
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simple contact representation height $=$ simple homotopy height
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visibility representation height
straight-line drawing height
inequalities are strict
gaps are nonconstant

## Pathwidth $\leq$ grid-major height

Pathwidth $=3$


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Grid-major height $\geq 2$ outerplanarity $-1 \geq n / 3-1$


Pathwidth $\leq$ grid-major height
Pathwidth $=3$
Grid-major height $\geq 2$ outerplanarity $-1 \geq n / 3-1$
$n / 6$ triangles will be nested, no matter the outer face


Outerplanarity $\leq$ grid-major height


Outerplanarity $\leq$ grid-major height
Grid-major height $\geq$ pathwidth $=\Omega(\log n)$


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Outerplanarity $\leq$ grid-major height
Grid-major height $\geq$ pathwidth $=\Omega(\log n)$ Outerplanarity $=2$


Nonsimple $\leq$ simple grid-major height


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Nonsimple $\leq$ simple grid-major height


Minor of and hence of $W \times 4$ grid

Nonsimple $\leq$ simple grid-major height


Minor of $\quad$ and hence of $W \times 4$ grid $\Rightarrow$ grid-major height $\leq 4$

Nonsimple $\leq$ simple grid-major height


Grid-major height $\leq 4$
Simple grid-major height $=\Omega(n)$ :

Nonsimple $\leq$ simple grid-major height


Grid-major height $\leq 4$
Simple grid-major height $=\Omega(n)$ :
Diameter of subgraph is $\Omega(n)$

Nonsimple $\leq$ simple grid-major height


Grid-major height $\leq 4$
Simple grid-major height $=\Omega(n)$ :
Diameter of subgraph is $\Omega(n)$
Some vertex in subgraph is far from 'outer face'

Nonsimple $\leq$ simple grid-major height


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## Nonsimple $\leq$ simple grid-major height



Grid-major height $\leq 4$
Simple grid-major height $=\Omega(n)$ :
Diameter of subgraph is $\Omega(n)$
Some vertex in subgraph is far from 'outer face'
That vertex splits some path in sweep in two pieces

## Nonsimple $\leq$ simple grid-major height



Grid-major height $\leq 4$
Simple grid-major height $=\Omega(n)$ :
Diameter of subgraph is $\Omega(n)$
Some vertex in subgraph is far from 'outer face'
That vertex splits some path in sweep in two pieces
At least one piece lies in subgraph, and is therefore long

## Simple grid-major height $\leq$ graph-drawing height

For series-parallel graphs, simple grid-major height is $O(\log n)$

Simple grid-major height $\leq$ graph-drawing height
For series-parallel graphs, simple grid-major height is $O(\log n)$

edge

series

parallel

## Simple grid-major height $\leq$ graph-drawing height

For series-parallel graphs, simple grid-major height is $O(\log n)$ Contact-representation with source/target in top/bottom-right

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parallel

Simple grid-major height $\leq$ graph-drawing height
For series-parallel graphs, simple grid-major height is $O(\log n)$
Contact-representation with source/target in top/bottom-right

series

parallel

Height increases (by 2 ) only if combined grids are similar height
$\Rightarrow$ grid-major height $=O(\log n)$

Simple grid-major height $\leq$ graph-drawing height
For series-parallel graphs, simple grid-major height is $O(\log n)$

There exist series-parallel graphs with
graph-drawing height $=\Omega\left(2^{\sqrt{\log n}}\right)$ [Frati10]

Simple grid-major height $\leq$ graph-drawing height
For series-parallel graphs, simple grid-major height is $O(\log n)$

There exist series-parallel graphs with
graph-drawing height $=\Omega\left(2^{\sqrt{\log n}}\right)$ [Frati10]
Triangulating them cannot decrease height

## Overview of results

2 outerplanarity -1 and pathwidth
grid-major height
contact representation height
$=$ homotopy height
simple grid-major height
simple contact representation height $=$ simple homotopy height
visibility representation height
straight-line drawing height inequalities are strict gaps are nonconstant

## Overview of results

Can we efficiently compute these parameters?
2 outerplanarity -1 and pathwidth (they are FPT in height)
grid-major height
contact representation height
$=$ homotopy height
simple contact representation height $=$ simple homotopy height
visibility representation height
straight-line drawing height
inequalities are strict
gaps are nonconstant

