

Chair for **INFORMATICS I** Efficient Algorithms and **Knowledge-Based Systems**



Drawing Planar Graphs with Few Segments on the Grid

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joint work with





Thomas Schneck Antonios Symvonis



of geometric entities in a drawing

of geometric entities in a drawing



of geometric entities in a drawing





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of geometric entities in a drawing







of geometric entities in a drawing



(strong) line cover number



of geometric entities in a drawing



(strong) line cover number



of geometric entities in a drawing





of geometric entities in a drawing





of geometric entities in a drawing





of geometric entities in a drawing





of geometric entities in a drawing





of geometric entities in a drawing





of geometric entities in a drawing



(strong) line cover number





segment number





of geometric entities in a drawing



(strong) line cover number





segment number





of geometric entities in a drawing



(strong) line cover number





segment number





of geometric entities in a drawing



(strong) line cover number





segment number





of geometric entities in a drawing



(strong) line cover number





segment number





of geometric entities in a drawing



(strong) line cover number





segment number





of geometric entities in a drawing



(strong) line cover number 7 6



segment number



arc number





of geometric entities in a drawing



(strong) line cover number 7 6



segment number



arc number





of geometric entities in a drawing



(strong) line cover number





segment number



arc number





of geometric entities in a drawing



(strong) line cover number





segment number



arc number




of geometric entities in a drawing



(strong) line cover number 7 6



segment number



arc number





(strong) line cover number

of geometric entities in a drawing

6

slope number





segment number

7



arc number





(strong) line cover number

of geometric entities in a drawing

6



slope number



segment number

7



arc number





(strong) line cover number

of geometric entities in a drawing

6



slope number



segment number

7



all other numbers are lower bounds

arc number





Class	Segments	
	Lower	Upper
	l	I

Class	Segments		
	Lower	Upper	
tree	θ/2 [1]	ϑ/2 [1]	
[1] Dujmović et a	1. 2007		

Class	Segments		
	Lower	Upper	
tree	θ/2 [1]	ϑ/2 [1]	
outerplanar	n [1]		
[1] Dujmović et a	1. 2007		

Class	Segments		
	Lower	Upper	r
tree	θ/2 [1]	ϑ/2	[1]
outerplanar	n [1]		
max. outerp.	n [1]	n	[1]
[1] Dujmović et a	1. 2007		

Class	Segments		
	Lower	Upper	
tree	θ/2 [1]	ϑ/2 [1]	
outerplanar	n [1]		
max. outerp.	n [1]	<i>n</i> [1]	
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]	
[1] Dujmović et a	1. 2007		

Class	Segments			
	Lowe	er	Upper	-
tree	ϑ/2 [[1]	ϑ/2	[1]
outerplanar	n [[1]		
max. outerp.	n [[1]	п	[1]
3-trees	2 <i>n</i> [[1]	2 <i>n</i>	[1]
2-connected	2 <i>n</i> [[1]		

Class	Segments			
	Lower	Upper		
tree	θ/2 [1]	ϑ/2 [1]		
outerplanar	n [1]			
max. outerp.	n [1]	<i>n</i> [1]		
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]		
2-connected	2 <i>n</i> [1]			
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]		

Class	Segments		
	Lower	Upper	
tree	ϑ/2 [1]	ϑ/2 [1]	
outerplanar	n [1]		
max. outerp.	n [1]	<i>n</i> [1]	
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]	
2-connected	2 <i>n</i> [1]		
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]	
cubic 3-conn.	n/2 [3]	<i>n</i> /2 [2]	

[1] Dujmović et al. 2007 [2] Igamberdiev et al. 2015 [3] Mondal et al. 2013

Class	Segments			
	Lower	Upper		
tree	θ/2 [1]	ϑ/2 [1]		
outerplanar	n [1]			
max. outerp.	n [1]	<i>n</i> [1]		
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]		
2-connected	2 <i>n</i> [1]			
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]		
cubic 3-conn.	n/2 [3]	<i>n</i> /2 [2]		
triangulation	2 <i>n</i> [4]	7 <i>n</i> /3 [4]		

Class	Segments		
	Lower	Upper	
tree	ϑ/2 [1]	ϑ/2 [1]	
outerplanar	n [1]		
max. outerp.	n [1]	<i>n</i> [1]	
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]	
2-connected	2 <i>n</i> [1]		
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]	
cubic 3-conn.	n/2 [3]	n/2 [2]	
triangulation	2 <i>n</i> [4]	7n/3 [4]	
4-conn. triang.	2 <i>n</i> [4]	9n/3 [4]	

Class	Segments		
	Lower	Upper	
tree	θ/2 [1]	ϑ/2 [1]	
outerplanar	n [1]		
max. outerp.	n [1]	n [1]	
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]	
2-connected	2 <i>n</i> [1]		
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]	
cubic 3-conn.	n/2 [3]	n/2 [2]	
triangulation	2 <i>n</i> [4]	7 <i>n</i> /3 [4]	
4-conn. triang.	2 <i>n</i> [4]	9n/3 [4]	
planar	2 <i>n</i> [4]	8 <i>n</i> /3 [4]	

Class	Segr	nents	Grid	Segments
	Lower	Upper	Segm.	Area
tree	θ/2 [1]	ϑ/2 [1]		
outerplanar	<i>n</i> [1]			
max. outerp.	n [1]	n [1]		
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]		
2-connected	2 <i>n</i> [1]			
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]		
cubic 3-conn.	n/2 [3]	<i>n</i> /2 [2]		
triangulation	2 <i>n</i> [4]	7 <i>n</i> /3 [4]		
4-conn. triang.	2 <i>n</i> [4]	9n/3 [4]		
planar	2 <i>n</i> [4]	8 <i>n</i> /3 [4]		

Class	Segments		Grid Segments	
	Lower	Upper	Segm.	Area
tree	ϑ/2 [1]	ϑ/2 [1]		
outerplanar	n [1]			
max. outerp.	n [1]	n [1]		
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]		
2-connected	2 <i>n</i> [1]			
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]		
cubic 3-conn.	n/2 [3]	n/2 [2]	n/2 [2]	$O(n) \times O(n)$
triangulation	2 <i>n</i> [4]	7 <i>n</i> /3 [4]		
4-conn. triang.	2 <i>n</i> [4]	9n/3 [4]		
planar	2 <i>n</i> [4]	8 <i>n</i> /3 [4]		

Class	Segments		Grid Segments	
	Lower	Upper	Segm.	Area
tree	θ/2 [1]	ϑ/2 [1]		
outerplanar	n [1]			
max. outerp.	n [1]	<i>n</i> [1]		
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]		
2-connected	2 <i>n</i> [1]			
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]		
cubic 3-conn.	n/2 [3]	<i>n</i> /2 [2]	n/2 [2]	$O(n) \times O(n)$
triangulation	2 <i>n</i> [4]	7 <i>n</i> /3 [4]	8n/3 [5]	$2^{O}(n\log n)$
4-conn. triang.	2 <i>n</i> [4]	9n/3 [4]		
planar	2 <i>n</i> [4]	8 <i>n</i> /3 [4]		

Class	Segments		Grid Segments	
	Lower	Upper	Segm.	Area
tree	θ/2 [1]	ϑ/2 [1]	3 <i>n</i> /4 [6]	$O(n^2) \times O(n^{1.58})$
outerplanar	n [1]			
max. outerp.	n [1]	n [1]		
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]		
2-connected	2 <i>n</i> [1]			
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]		
cubic 3-conn.	n/2 [3]	<i>n</i> /2 [2]	n/2 [2]	$O(n) \times O(n)$
triangulation	2 <i>n</i> [4]	7 <i>n</i> /3 [4]	8n/3 [5]	$2^{O}(n\log n)$
4-conn. triang.	2 <i>n</i> [4]	9n/3 [4]		
planar	2 <i>n</i> [4]	8 <i>n</i> /3 [4]		

Class	Segments		Grid Segments	
	Lower	Upper	Segm.	Area
tree	θ/2 [1]	ϑ/2 [1]	3n/4 [6] ϑ/2 [6]	$O(n^2) \times O(n^{1.58})$ quasipolynomial
outerplanar	n [1]			
max. outerp.	n [1]	n [1]		
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]		
2-connected	2 <i>n</i> [1]			
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]		
cubic 3-conn.	n/2 [3]	<i>n</i> /2 [2]	n/2 [2]	$O(n) \times O(n)$
triangulation	2 <i>n</i> [4]	7 <i>n</i> /3 [4]	8n/3 [5]	$2^{O}(n\log n)$
4-conn. triang.	2 <i>n</i> [4]	9n/3 [4]		
planar	2 <i>n</i> [4]	8 <i>n</i> /3 [4]		

Class	Segments		Grid Segments	
	Lower	Upper	Segm.	Area
tree	θ/2 [1]	θ/2 [1]	3 <i>n</i> /4 [6] <i>v</i> /2 [6]	$\begin{array}{ c }O(n^2) \times O(n^{1.58})\\ \text{quasipolynomial}\end{array}$
outerplanar	n [1]			
max. outerp.	n [1]	n [1]	3 <i>n</i> /2 [6]	$O(n) \times O(n^2)$
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]		
2-connected	2 <i>n</i> [1]			
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]		
cubic 3-conn.	n/2 [3]	n/2 [2]	n/2 [2]	$O(n) \times O(n)$
triangulation	2 <i>n</i> [4]	7 <i>n</i> /3 [4]	8n/3 [5]	$2^{O}(n\log n)$
4-conn. triang.	2 <i>n</i> [4]	9n/3 [4]		
planar	2 <i>n</i> [4]	8 <i>n</i> /3 [4]		

Class	Segments		Grid Segments	
	Lower	Upper	Segm.	Area
tree	θ/2 [1]	θ/2 [1]	3n/4 [6] v/2 [6]	$\begin{array}{ c } O(n^2) \times O(n^{1.58}) \\ \text{quasipolynomial} \end{array}$
outerplanar	n [1]			
max. outerp.	n [1]	<i>n</i> [1]	3n/2 [6]	$O(n) \times O(n^2)$
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]	8n/3 [6]	$O(n) \times O(n^2)$
2-connected	2 <i>n</i> [1]			
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]		
cubic 3-conn.	n/2 [3]	n/2 [2]	n/2 [2]	$O(n) \times O(n)$
triangulation	2 <i>n</i> [4]	7 <i>n</i> /3 [4]	8n/3 [5]	$2^{O}(n\log n)$
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Class	Segments		Grid Segments	
	Lower	Upper	Segm.	Area
tree	θ/2 [1]	ϑ/2 [1]	3n/4 [6] ϑ/2 [6]	$O(n^2) \times O(n^{1.58})$ quasipolynomial
outerplanar	n [1]			
max. outerp.	n [1]	n [1]	3n/2 [6]	$O(n) \times O(n^2)$
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]	8n/3 [6]	$O(n) \times O(n^2)$
2-connected	2 <i>n</i> [1]			
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cubic 3-conn.	n/2 [3]	n/2 [2]	n/2 [2]	$O(n) \times O(n)$
triangulation	2 <i>n</i> [4]	7 <i>n</i> /3 [4]	8n/3 [5]	$2^{O}(n\log n)$
4-conn. triang.	2 <i>n</i> [4]	9n/3 [4]		
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Class	Segments		Grid Segments	
	Lower	Upper	Segm.	Area
tree	θ/2 [1]	ϑ/2 [1]	3n/4 [6] ϑ/2 [6]	$O(n^2) \times O(n^{1.58})$ guasipolynomial
outerplanar	n [1]		,	
max. outerp.	<i>n</i> [1]	n [1]	3n/2 [6]	$O(n) \times O(n^2)$
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]	8n/3 [6]	$O(n) \times O(n^2)$
2-connected	2 <i>n</i> [1]			
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cubic 3-conn.	n/2 [3]	n/2 [2]	n/2 [2]	$O(n) \times O(n)$
triangulation	2 <i>n</i> [4]	7 <i>n</i> /3 [4]	8n/3 [5]	$2^{O}(n\log n)$
4-conn. triang.	2 <i>n</i> [4]	9n/3 [4]		
planar	2 <i>n</i> [4]	8 <i>n</i> /3 [4]		

Tree Drawings Tree T *n* vtcs D О





Tree *T n* vtcs



Remove β deg-2 vtcs

Tree *T n* vtcs

 $\Rightarrow \text{Tree } T'$ $n - \beta \text{ vtcs}$



Remove β deg-2 vtcs

Tree *T n* vtcs

 $\Rightarrow \text{Tree } T'$ $n - \beta \text{ vtcs}$



Remove β deg-2 vtcs

 α leaves

Tree *T n* vtcs

 $\Rightarrow \text{Tree } T'$ $n - \beta \text{ vtcs}$



Remove β deg-2 vtcs

 α leaves

Tree *T n* vtcs

 $\Rightarrow \text{Tree } T'$ $n - \beta \text{ vtcs}$



Remove β deg-2 vtcs

 α leaves

Tree *T n* vtcs

 $\Rightarrow \text{Tree } T'$ $n - \beta \text{ vtcs}$



Remove β deg-2 vtcs

Remove α leaves

Tree *T n* vtcs

 $\Rightarrow \text{Tree } T'$ $n - \beta \text{ vtcs}$



 $\Rightarrow \text{Tree } T'' \\ n - \alpha - \beta \text{ vtcs}$





Remove β deg-2 vtcs

Remove α leaves

Tree *T n* vtcs

 $\Rightarrow \text{Tree } T'$ $n - \beta \text{ vtcs}$



$$\Rightarrow \text{Tree } T'' \\ n - \alpha - \beta \text{ vtcs}$$





Remove β deg-2 vtcs

Remove α leaves



 $n - \alpha - \beta$ segment

Tree *T n* vtcs

 $\Rightarrow \text{Tree } T'$ $n - \beta \text{ vtcs}$



 $\Rightarrow \text{Tree } T'' \\ n - \alpha - \beta \text{ vtcs}$





Remove β deg-2 vtcs

Remove α leaves + $\alpha/2$ segments



 $n - \alpha - \beta$ segment
Tree *T n* vtcs

 $\Rightarrow \text{Tree } T'$ $n - \beta \text{ vtcs}$



 $\Rightarrow \text{Tree } T'' \\ n - \alpha - \beta \text{ vtcs}$





Remove β deg-2 vtcs

Remove α leaves

+ $\alpha/2$ segments





Tree *T n* vtcs

 $\Rightarrow \text{Tree } T'$ $n - \beta \text{ vtcs}$



 $\Rightarrow \text{Tree } T'' \\ n - \alpha - \beta \text{ vtcs}$





Remove β deg-2 vtcs + 0 segments

Remove α leaves + $\alpha/2$ segments





Tree *T n* vtcs

 $\Rightarrow \text{Tree } T'$ $n - \beta \text{ vtcs}$



 $\Rightarrow \text{Tree } T'' \\ n - \alpha - \beta \text{ vtcs}$





Remove β deg-2 vtcs

+ 0 segments



Remove α leaves + $\alpha/2$ segments



segments



Tree *T n* vtcs

 $\Rightarrow \text{Tree } T'$ $n - \beta \text{ vtcs}$



 $\Rightarrow \text{Tree } T'' \\ n - \alpha - \beta \text{ vtcs}$





Remove β deg-2 vtcs

+ 0 segments



Remove α leaves + $\alpha/2$ segments





 $n - \alpha - \beta$ segment

Tree *T n* vtcs

 $\Rightarrow \text{Tree } T'$ $n - \beta \text{ vtcs}$



 $\Rightarrow \text{Tree } T'' \\ n - \alpha - \beta \text{ vtcs}$



 $\alpha > (n-\beta)/2$

Remove β deg-2 vtcs

+ 0 segments



Remove α leaves + $\alpha/2$ segments

 $\frac{1}{n-\alpha/2-\beta}$

segments



Tree *T n* vtcs



 $\Rightarrow \text{Tree } T'$ $n - \beta \text{ vtcs}$



 $\Rightarrow \text{Tree } T'' \\ n - \alpha - \beta \text{ vtcs}$



 $\alpha > (n-\beta)/2$

Remove β deg-2 vtcs

+ 0 segments



Remove α leaves + $\alpha/2$ segments





Tree Drawings





Tree Drawings





 v_{o}















































 v_{o}























(2) Layout $v + \overset{\wedge}{\Delta} \overset{\wedge}{\Delta} \overset{\wedge}{\Delta} \overset{\wedge}{\Delta}$







- (1) Draw $\mathring{\Delta} \mathring{\Delta} \mathring{\Delta} \mathring{\Delta}$
- (2) Layout $v + \Delta \Delta \Delta \Delta$
- (3) Add 🛛







- (1) Draw $\mathring{\Delta} \mathring{\Delta} \mathring{\Delta} \mathring{\Delta}$
- (2) Layout $v + \overset{\wedge}{\Delta} \overset{\wedge}{\Delta} \overset{\wedge}{\Delta} \overset{\wedge}{\Delta}$
- (3) Add 🛛







- (1) Draw $\mathring{\Delta} \mathring{\Delta} \mathring{\Delta} \mathring{\Delta}$
- (2) Layout $v + \overset{\wedge}{\Delta} \overset{\wedge}{\Delta} \overset{\wedge}{\Delta} \overset{\wedge}{\Delta}$
- (3) Add 🛛







- (1) Draw $\mathring{\Delta} \mathring{\Delta} \mathring{\Delta} \mathring{\Delta}$
- (2) Layout $v + \overset{\wedge}{\Delta} \overset{\wedge}{\Delta} \overset{\wedge}{\Delta} \overset{\wedge}{\Delta}$
- (3) Add 🛛





- Tree Drawings
- (1) Draw $\mathring{\Delta} \mathring{\Delta} \mathring{\Delta} \mathring{\Delta}$
- (2) Layout $v + \overset{\wedge}{\square} \overset{\wedge}{\square} \overset{\wedge}{\square} \overset{\wedge}{\square}$
- (3) Add 🛛
- (4) Sort by #





- Tree Drawings
- (1) Draw $\mathring{\Delta} \mathring{\Delta} \mathring{\Delta} \mathring{\Delta}$
- (2) Layout $v + \overset{\wedge}{\square} \overset{\wedge}{\square} \overset{\wedge}{\square} \overset{\wedge}{\square}$
- (3) Add 🛛
- (4) Sort by #







- (4) Sort by #
- (5) Place + on common segments in order





- (4) Sort by #
- (5) Place + on common segments in order





- (4) Sort by #
- (5) Place + on common segments in order





- (4) Sort by #
- (5) Place + on common segments in order





- (4) Sort by #
- (5) Place + on common segments in order





- (4) Sort by #
- (5) Place + on common segments in order




- (4) Sort by #
- (5) Place + on common segments in order





- (4) Sort by #
- (5) Place + □ on common segments in order































Improved Results

Class	Segments		Grid Segments	
	Lower	Upper	Segm.	Area
tree	θ/2 [1]	ϑ/2 [1]	$\frac{3n}{4}$ [6] $\frac{9}{2}$ [6]	$O(n^2) \times O(n^{1.58})$ quasipolynomial
outerplanar	n [1]			
max. outerp.	n [1]	n [1]	3n/2 [6]	$O(n) \times O(n^2)$
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]	8n/3 [6]	$O(n) \times O(n^2)$
2-connected	2 <i>n</i> [1]			
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]		
cubic 3-conn.	n/2 [3]	<i>n</i> /2 [2]	n/2 [2]	$O(n) \times O(n)$
triangulation	2 <i>n</i> [4]	7 <i>n</i> /3 [4]	8n/3 [5]	$2^{O}(n\log n)$
4-conn. triang.	2 <i>n</i> [4]	9n/3 [4]		
planar	2 <i>n</i> [4]	8 <i>n</i> /3 [4]		

[1] Dujmović et al. 2007 [2] Igamberdiev et al. 2015 [3] Mondal et al. 2013
[4] Durocher & Mondal 2014 [5] Mondal 2016 [6] Hültenschmidt et al. 2017

Improved Results

Class	Segments		Grid Segments	
	Lower	Upper	Segm.	Area
tree	ϑ/2 [1]	$\vartheta/2$ [1]	3n/4	$n \times n$
outerplanar	n [1]	, , ,	V /2 [6]	quasipolynomial
max. outerp.	n [1]	<i>n</i> [1]	3n/2 [6]	$O(n) \times O(n^2)$
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]	8n/3 [6]	$O(n) \times O(n^2)$
2-connected	2 <i>n</i> [1]			
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]		
cubic 3-conn.	n/2 [3]	<i>n</i> /2 [2]	n/2 [2]	$O(n) \times O(n)$
triangulation	2 <i>n</i> [4]	7 <i>n</i> /3 [4]	8n/3 [5]	$2^{O}(n\log n)$
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[4] Durocher & Mondal 2014 [5] Mondal 2016 [6] Hültenschmidt et al. 2017













[Chiang, Lin, Lu '05]





[Chiang, Lin, Lu '05]





neighbors of *v* in circ. order: (1) parent

[Chiang, Lin, Lu '05]





- (1) parent
- (2) $N^+(v)$: diff. subtree (left)

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[Angelini et al. '12]

Assign angle interval to each vtx



Slope-Disjoint Drawing of a Tree [Angelini et al. '12]

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Slope-disjoint drawing of orderly spanning tree on $O(n) \times O(n^2)$ grid \Rightarrow planar (monotone) drawing on $O(n) \times O(n^2)$ grid

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doesn't change the slopes!

















ccw pre-order traversal reuse slope whenever possible





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highest slope: *n*





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highest slope: *n* max. width: *n*





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$$\Rightarrow n \times n^2$$
 grid





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 \Rightarrow *n* × *n*² grid , 1 segment per leaf



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[Miura, Azuma, Nishizeki '05]

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*T*₁, *T*₂, *T*₃ Schnyder realizer of 3-conn. planar graph \Rightarrow ≤ 2*n* + 1 leaves in total in *T*₁, *T*₂, *T*₃

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3-conn. planar graph $\Rightarrow (8n - 14)/3$ segments, $O(n) \times O(n^2)$ grid

Class	Segments		Grid Segments	
	Lower	Upper	Segm.	Area
tree	θ/2 [1]	ϑ/2 [1]	3n/4 v/2 [6]	$n \times n$ quasipolynomial
outerplanar	n [1]			
max. outerp.	n [1]	n [1]	3 <i>n</i> /2 [6]	$O(n) \times O(n^2)$
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]	8n/3 [6]	$O(n) \times O(n^2)$
2-connected	2 <i>n</i> [1]			
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]		
cubic 3-conn.	n/2 [3]	n/2 [2]	n/2 [2]	$O(n) \times O(n)$
triangulation	2 <i>n</i> [4]	7 <i>n</i> /3 [4]	8n/3 [5]	$2^{O}(n\log n)$
4-conn. triang.	2 <i>n</i> [4]	9n/3 [4]		
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cubic 3-conn.	n/2 [3]	n/2 [2]	n/2 [2]	$O(n) \times O(n)$
triangulation	2 <i>n</i> [4]	7n/3 [4]	8n/3 [5]	$2^{O}(n\log n)$
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4-conn. triang.	2 <i>n</i> [4]	9n/3 [4]	5 <i>n</i> /2	$O(n) \times O(n^2)$
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planar	2 <i>n</i> [4]	8 <i>n</i> /3 [4]	17n/6	$O(n) \times O(n^2)$
New Results

Class	Segments		Grid Segments	
	Lower	Upper	Segm.	Area
tree	ϑ/2 [1]	ϑ/2 [1]	3n/4 ϑ/2 [6]	$n \times n$ quasipolynomial
outerplanar	n [1]		7n/4	$O(n) \times O(n^2)$
max. outerp.	n [1]	n [1]	3n/2 [6]	$O(n) \times O(n^2)$
3-trees	2 <i>n</i> [1]	2 <i>n</i> [1]	8n/3 [6]	$O(n) \times O(n^2)$
2-connected	2 <i>n</i> [1]		17n/6	$O(n) \times O(n^2)$
3-connected	2 <i>n</i> [1]	5 <i>n</i> /2 [1]	8n/3	$O(n) \times O(n^2)$
cubic 3-conn.	n/2 [3]	n/2 [2]	n/2 [2]	$O(n) \times O(n)$
triangulation	2 <i>n</i> [4]	7 <i>n</i> /3 [4]	8n/3 [5]	$2^{O}(n\log n)$
4-conn. triang.	2 <i>n</i> [4]	9n/3 [4]	5 <i>n</i> /2	$O(n) \times O(n^2)$
planar	2 <i>n</i> [4]	8 <i>n</i> /3 [4]	17n/6	$O(n) \times O(n^2)$

[1] Dujmović et al. 2007 [2] Igamberdiev et al. 2015 [3] Mondal et al. 2013
[4] Durocher & Mondal 2014 [5] Mondal 2016 [6] Hültenschmidt et al. 2017